

Parallel Graph Algorithms on the Xeon Phi Coprocessor

Master Thesis presentation

Dennis Felsing | 2015-09-07

INSTITUTE OF THEORETICAL INFORMATICS, RESEARCH GROUP PARALLEL COMPUTING





Complex Network: graph with non-trivial topology

Occur in social networks, cell biology, the internet...



Map of the Internet, http://www.opte.org/maps/



- Complex Network: graph with non-trivial topology
- Occur in social networks, cell biology, the internet...
- We consider two existing algorithms:



- Complex Network: graph with non-trivial topology
- Occur in social networks, cell biology, the internet...
- We consider two existing algorithms:
- Graph Generation
 - Create realistic complex networks with generator and parameters
 - Preserves privacy and confidentiality
 - No need to transfer big data
 - Scale to smaller and bigger graphs





- Complex Network: graph with non-trivial topology
- Occur in social networks, cell biology, the internet...
- We consider two existing algorithms:
- Graph Generation
 - Create realistic complex networks with generator and parameters
 - Preserves privacy and confidentiality
 - No need to transfer big data
 - Scale to smaller and bigger graphs
- Graph Drawing
 - Lay out graph visually, mainly for human perception



- Data sets grow fast: Internet size doubles every 5 years
- Clock frequency of processors stagnated in last decade





- Data sets grow fast: Internet size doubles every 5 years
- Clock frequency of processors stagnated in last decade
- Instead more parallelism in processors



- Data sets grow fast: Internet size doubles every 5 years
- Clock frequency of processors stagnated in last decade
- Instead more parallelism in processors
- Modern GPUs as even more parallel alternative:
 - Massively parallel, thousands of cores
 - Large performance increases with more parallelization



- Data sets grow fast: Internet size doubles every 5 years
- Clock frequency of processors stagnated in last decade
- Instead more parallelism in processors
- Modern GPUs as even more parallel alternative:
 - Massively parallel, thousands of cores
 - Large performance increases with more parallelization
 - General purpose programming more difficult
 - Graph algorithms with irregular data access challenging



- Data sets grow fast: Internet size doubles every 5 years
- Clock frequency of processors stagnated in last decade
- Instead more parallelism in processors
- Modern GPUs as even more parallel alternative:
 - Massively parallel, thousands of cores
 - Large performance increases with more parallelization
 - General purpose programming more difficult
 - Graph algorithms with irregular data access challenging
- Intel Xeon Phi as alternative to the alternative:
 - 60 cores, more than a CPU, fewer than a GPU
 - Similar to program as CPU
 - Viable choice for graph algorithms?





- Data sets grow fast: Internet size doubles every 5 years
- Clock frequency of processors stagnated in last decade
- Instead more parallelism in processors
- Modern GPUs as even more parallel alternative:
 - Massively parallel, thousands of cores
 - Large performance increases with more parallelization
 - General purpose programming more difficult
 - Graph algorithms with irregular data access challenging
- Intel Xeon Phi as alternative to the alternative:
 - 60 cores, more than a CPU, fewer than a GPU
 - Similar to program as CPU
 - Viable choice for graph algorithms?
- \Rightarrow Port two graph algorithms to Xeon Phi and evaluate the porting and their performance



Overview





Xeon Phi Coprocessor

- Hardware Architecture
- Programming
- Generation of Massive Complex Networks
 - Algorithm
 - Results



- Algorithm
- Results



- Xeon Phi 5110P used, 60 in-order cores at 1 GHz
- Simple cores based on original Pentium design from 1994
- Augmented with 64-bit support (not x86-64)





- Xeon Phi 5110P used, 60 in-order cores at 1 GHz
- Simple cores based on original Pentium design from 1994
- Augmented with 64-bit support (not x86-64)

| | Host System | Accelerator Card |
|-------------------|--------------------------|------------------|
| Name | $2 \times$ Xeon E5-2680 | Xeon Phi 5110P |
| Release Date | 2012 | 2012 |
| Clock Frequency | 2.7 GHz | 1.05 GHz |
| Cores | 2 $	imes$ 8 (32 threads) | 60 (240 threads) |
| RAM Capacity | 256 GB | 8 GB |
| RAM Bandwidth | 51 GB/s | 320 GB/s |
| SIMD Instructions | MMX, SSE, AVX (256 bit) | IMCI (512 bit) |



- Xeon Phi 5110P used, 60 in-order cores at 1 GHz
- Simple cores based on original Pentium design from 1994
- Augmented with 64-bit support (not x86-64)

| | Host System | Accelerator Card |
|-------------------|--------------------------|------------------|
| Name | 2× Xeon E5-2680 | Xeon Phi 5110P |
| Release Date | 2012 | 2012 |
| Clock Frequency | 2.7 GHz | 1.05 GHz |
| Cores | 2 $	imes$ 8 (32 threads) | 60 (240 threads) |
| RAM Capacity | 256 GB | 8 GB |
| RAM Bandwidth | 51 GB/s | 320 GB/s |
| SIMD Instructions | MMX, SSE, AVX (256 bit) | IMCI (512 bit) |

 \Rightarrow Parallelization and vectorization necessary to reach high performance



- Xeon Phi 5110P used, 60 in-order cores at 1 GHz
- Simple cores based on original Pentium design from 1994
- Augmented with 64-bit support (not x86-64)





- Presented as a regular computer, stripped down Linux, SSH
- Parallelization (Multiple Instruction Multiple Data) methods:
 - OpenMP
 - OpenMP Offloading
 - Cilk Plus
 - Threading Building Blocks
 - MPI
- Vectorization (Single Instruction Multiple Data) methods



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization (Multiple Instruction Multiple Data) methods:
 - OpenMP

```
float a[MAX], b[MAX], c[MAX];
for (i = 0; i < MAX; i++)
   c[i] = a[i] + b[i];</pre>
```

- OpenMP Offloading
- Cilk Plus
- Threading Building Blocks
- MPI
- Vectorization (Single Instruction Multiple Data) methods



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization (Multiple Instruction Multiple Data) methods:
 - OpenMP

```
float a[MAX], b[MAX], c[MAX];
#pragma omp parallel for
for (i = 0; i < MAX; i++)
   c[i] = a[i] + b[i];</pre>
```

- OpenMP Offloading
- Cilk Plus
- Threading Building Blocks
- MPI
- Vectorization (Single Instruction Multiple Data) methods



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization (Multiple Instruction Multiple Data) methods:
 - OpenMP
 - OpenMP Offloading

```
float a[MAX], b[MAX], c[MAX];
#pragma offload target(mic) in(a, b) out(c)
{
    #pragma omp parallel for
    for (i = 0; i < MAX; i++)
        c[i] = a[i] + b[i];
}
Cilk Plus
Threading Building Blocks
MPI</pre>
```



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization (Multiple Instruction Multiple Data) methods:
 - OpenMP
 - OpenMP Offloading

```
float a[MAX], b[MAX], c[MAX];
#pragma offload target(mic:0) in(a, b) out(c) signal(c)
{
    #pragma omp parallel for
    for (i = 0; i < MAX; i++)
        c[i] = a[i] + b[i];
}
#pragma offload_wait target(mic:0) wait(c)
Cilk Plus
The W D D D d
</pre>
```

- Threading Building Blocks
- MPI



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization (Multiple Instruction Multiple Data) methods
- Vectorization (Single Instruction Multiple Data) methods:
 - Manual Vectorization: c = _mm512_add_ps(a, b)





- Auto-Vectorization
- Cilk Plus



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization (Multiple Instruction Multiple Data) methods
- Vectorization (Single Instruction Multiple Data) methods:
 - Manual Vectorization: c = _mm512_add_ps(a, b)





- Auto-Vectorization
- Cilk Plus



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization (Multiple Instruction Multiple Data) methods
- Vectorization (Single Instruction Multiple Data) methods:
 - Manual Vectorization: c = _mm512_add_ps(a, b)





- Auto-Vectorization
- Cilk Plus



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization (Multiple Instruction Multiple Data) methods
- Vectorization (Single Instruction Multiple Data) methods:
 - Manual Vectorization: c = _mm512_add_ps(a, b)





- Auto-Vectorization
- Cilk Plus

Dennis Felsing - Parallel Graph Algorithms on the Xeon Phi Coprocessor

512 bits



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization (Multiple Instruction Multiple Data) methods
- Vectorization (Single Instruction Multiple Data) methods:
 - Manual Vectorization: c = _mm512_add_ps(a, b)

a 0123456789012345012345 32 bits (float) +

b 98765432109876549876543210987654

- Auto-Vectorization
- Cilk Plus



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization (Multiple Instruction Multiple Data) methods
- Vectorization (Single Instruction Multiple Data) methods:
 - Manual Vectorization
 - Auto-Vectorization

```
float a[MAX] __attribute__((aligned(64)));
float b[MAX] __attribute__((aligned(64)));
float c[MAX] __attribute__((aligned(64)));
#pragma simd
for (i = 0; i < MAX; i++)
    c[i] = a[i] + b[i];
or D</pre>
```

Cilk Plus



Desired properties of a complex network:

- Scale-Free: No typical vertex degree
 - \Rightarrow Degree distribution follows power law





Desired properties of a complex network:

- Scale-Free: No typical vertex degree
 - \Rightarrow Degree distribution follows power law
- Small-World: All nodes connected by short paths





Desired properties of a complex network:

- Scale-Free: No typical vertex degree
 - \Rightarrow Degree distribution follows power law
- Small-World: All nodes connected by short paths
- \Rightarrow Generator using hyperbolic geometry performs well in both properties





 Exponential expansion of space in hyperbolic geometry: Area of circle grows exponentially with distance from center
 Natural embedding of graphs with tree-like structure
 May also be good for generating graphs



M.C. Escher: Circle Limit IV

Dennis Felsing - Parallel Graph Algorithms on the Xeon Phi Coprocessor



- Exponential expansion of space in hyperbolic geometry: Area of circle grows exponentially with distance from center
 Natural embedding of graphs with tree-like structure
 May also be good for generating graphs
- Hyperbolic generator: Distribute vertices in hyperbolic plane
 Edge when two vertices are close to each other (hyperbolic circle)





- Exponential expansion of space in hyperbolic geometry: Area of circle grows exponentially with distance from center
 Natural embedding of graphs with tree-like structure
 May also be good for generating graphs
- Hyperbolic generator: Distribute vertices in hyperbolic plane
 Edge when two vertices are close to each other (hyperbolic circle)
- Poincaré disk model: Mapping to Euclidean unit disk
 - \Rightarrow Neighborhood transformed to Euclidean circle





- Exponential expansion of space in hyperbolic geometry: Area of circle grows exponentially with distance from center
 Natural embedding of graphs with tree-like structure
 May also be good for generating graphs
- Hyperbolic generator: Distribute vertices in hyperbolic plane
 Edge when two vertices are close to each other (hyperbolic circle)
- Poincaré disk model: Mapping to Euclidean unit disk ⇒ Neighborhood transformed to Euclidean circle
- Polar quadtree: Efficiently determine neighborhood





- Exponential expansion of space in hyperbolic geometry: Area of circle grows exponentially with distance from center
 Natural embedding of graphs with tree-like structure
 May also be good for generating graphs
- Hyperbolic generator: Distribute vertices in hyperbolic plane
 Edge when two vertices are close to each other (hyperbolic circle)
- Poincaré disk model: Mapping to Euclidean unit disk ⇒ Neighborhood transformed to Euclidean circle
- Polar quadtree: Efficiently determine neighborhood
 - \Rightarrow Subquadratic running time $O((n^{3/2} + m) \log n)$



Network Generation: Implementation



- Implementation part of NetworKit, high level C++11 code
- NetworKit ported to Intel C++ compiler 15.0 and Xeon Phi
- Working around compiler restrictions and bugs with *constexprs*, implicit conversions, null pointers, the standard library, and function traits on lambdas

 Three execution modes implemented and tested: No offloading: Entire code runs on Xeon Phi, not enough memory Full offloading: Offload parts of the calculation, keep results in memory of host system
 Partial offloading: Offload part of the calculation, other part on the host system



- Many memory allocations during algorithm to create dynamically sized lists of neighbors
- Allocations (malloc) are locking in default C library glibc
- On Xeon Phi more threads run in parallel than on CPU, so more allocations block each other



- Many memory allocations during algorithm to create dynamically sized lists of neighbors
- Allocations (malloc) are locking in default C library glibc
- On Xeon Phi more threads run in parallel than on CPU, so more allocations block each other
- \Rightarrow Use non-locking allocations of Intel's Threading Building Blocks





- Many memory allocations during algorithm to create dynamically sized lists of neighbors
- Allocations (malloc) are locking in default C library glibc
- On Xeon Phi more threads run in parallel than on CPU, so more allocations block each other
- \Rightarrow Use non-locking allocations of Intel's Threading Building Blocks
- \Rightarrow Reduce number of (re)allocations by reusing memory and preallocating expected size





 Tuning parameters of the algorithm: Capacity: Maximum number of vertices in a leaf cell before split Balance: Share of area in outer children when splitting





 Tuning parameters of the algorithm: Capacity: Maximum number of vertices in a leaf cell before split Balance: Share of area in outer children when splitting



 \Rightarrow Imbalanced quadtree with greater space to outer children



■ Transferring parts of graph back to host system is slow ⇒ Double Buffering on Phi and host

Buffering for full offloading:





Transferring parts of graph back to host system is slow
 Double Buffering on Phi and host



Graph Drawing: Algorithm



We assume to have graphs with predefined target edge lengths

- Full Stress Model: Physical springs connecting all pairs of vertices
- Maxent-Stress Model: Minimize stress, maximize entropy:



Graph Drawing: Algorithm



We assume to have graphs with predefined target edge lengths

- Full Stress Model: Physical springs connecting all pairs of vertices
- Maxent-Stress Model: Minimize stress, maximize entropy:



Multilevel Maxent-Stress Algorithm:

- Minimize maxent-stress by clustering graph in multiple levels of hierarchy
- Contract clusters into new supervertices
- Iteratively solve maxent-stress on each finer level
- \Rightarrow Parallelizes well

Graph Drawing: Results



- Preliminary implementation of multilevel maxent-stress graph drawing algorithm
- Parallelized with OpenMP
- Graphs of interest (< 10⁷ edges) easily fit into Xeon Phi memory ⇒ No expensive offloading necessary
- Source code libraries had to be fixed for ICPC
- No major dynamic allocations in algorithm ⇒ Intel TBB's *malloc* has small effect
- Inner loop vectorizes well when isolated: speedup factor 7.0 Smaller effect when embedded in real program

 Other calculations at same time, hyper-threading, memory connection busy

Graph Drawing: Results



| Graph | п | т | Description | Phi | Host |
|---------------|---------|---------|------------------------|-------|--------|
| nyc | 264 346 | 365 050 | Road Network | 960.0 | 1845.9 |
| luxembourg | 114 599 | 119666 | Road Network | 89.5 | 166.5 |
| commanche | 7920 | 11 880 | Helicopter Mesh | 2.6 | 3.5 |
| rajat06 | 10 922 | 18061 | Circuit Simulation | 3.5 | 4.5 |
| delaunay_n15 | 32 768 | 98 274 | Delaunay Triangulation | 7.8 | 9.0 |
| rgg_n_2_15_s0 | 32 768 | 160 240 | Random Graph | 5.3 | 3.6 |



- All graphs small enough to fit into Xeon Phi memory
 - \Rightarrow Executed on Xeon Phi directly
- Good speedup for large sparse graphs

Conclusion and Outlook



- Complex algorithms and their framework/libraries ported to Xeon Phi
- Good scaling in both algorithms on Xeon Phi
- Offloading large amounts of data too expensive
- Graph drawing algorithm outperforms two-socket Intel Xeon system, especially on sparse graphs
- Future Research: Direct comparison between graph algorithms on GPU and Xeon Phi
- New Xeon Phi "Knights Landing" this year with modern cores and 384 GB of memory

[Appendix] Xeon Phi: Layout Overview





Dennis Felsing - Parallel Graph Algorithms on the Xeon Phi Coprocessor

17/16

[Appendix] Xeon Phi: Core Pipeline







- Presented as a regular computer, stripped down Linux, SSH
- Parallelization methods:
 - OpenMP
 - OpenMP Offloading
 - Cilk Plus
 - Threading Building Blocks
 - MPI
- Vectorization (Single Instruction Multiple Data) methods



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization methods:
 - OpenMP

float a[MAX], b[MAX], c[MAX];

```
for (i = 0; i < MAX; i++)
c[i] = a[i] + b[i];</pre>
```

- OpenMP Offloading
- Cilk Plus
- Threading Building Blocks
- MPI
- Vectorization (Single Instruction Multiple Data) methods



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization methods:
 - OpenMP

float a[MAX], b[MAX], c[MAX];
#pragma omp parallel for
for (i = 0; i < MAX; i++)
 c[i] = a[i] + b[i];</pre>

- OpenMP Offloading
- Cilk Plus
- Threading Building Blocks
- MPI



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization methods:
 - OpenMP
 - OpenMP Offloading

```
float a[MAX], b[MAX], c[MAX];
#pragma offload target(mic) in(a, b) out(c)
{
    #pragma omp parallel for
    for (i = 0; i < MAX; i++)
        c[i] = a[i] + b[i];
}
Cilk Plus
Threading Building Blocks
MPI</pre>
```



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization methods:
 - OpenMP
 - OpenMP Offloading

```
float a[MAX], b[MAX], c[MAX];
#pragma offload target(mic:0) in(a, b) out(c) signal(c)
{
    #pragma omp parallel for
    for (i = 0; i < MAX; i++)
        c[i] = a[i] + b[i];
}
#pragma offload_wait target(mic:0) wait(c)
Cilk Plus
Thereadian Duridian Diadua</pre>
```

- Threading Building Blocks
- MPI



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization methods:
 - OpenMP
 - OpenMP Offloading
 - Cilk Plus

```
float a[MAX], b[MAX], c[MAX];
cilk_for (i = 0; i < MAX; i++)
c[i] = a[i] + b[i];
```

Threading Building Blocks

MPI



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization methods:
 - OpenMP
 - OpenMP Offloading
 - Cilk Plus
 - Threading Building Blocks

```
float a[MAX], b[MAX], c[MAX];
parallel_for(size_t(0), MAX, size_t(1), [=](size_t i) {
    c[i] = a[i] + b[i];
});
MPI
```



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization methods
- Vectorization (Single Instruction Multiple Data) methods:
 - Manual Vectorization: c = _mm512_add_ps(a, b)
 - 1 #include <immintrin.h>
 - 2 float a[MAX] __attribute__((aligned(64)));
 - 3 float b[MAX] __attribute__((aligned(64)));
 - 4 float c[MAX] __attribute__((aligned(64)));

8

1

Auto-Vectorization

```
Cilk Plus
```



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization methods
- Vectorization (Single Instruction Multiple Data) methods:
 - Manual Vectorization
 - Auto-Vectorization

```
float a[MAX] __attribute__((aligned(64)));
float b[MAX] __attribute__((aligned(64)));
float c[MAX] __attribute__((aligned(64)));
#pragma simd
for (i = 0; i < MAX; i++)
    c[i] = a[i] + b[i];
even be</pre>
```

Cilk Plus



- Presented as a regular computer, stripped down Linux, SSH
- Parallelization methods
- Vectorization (Single Instruction Multiple Data) methods:
 - Manual Vectorization
 - Auto-Vectorization
 - Cilk Plus

```
float a[MAX] __attribute__((aligned(64)));
float b[MAX] __attribute__((aligned(64)));
float c[MAX] __attribute__((aligned(64)));
c[i:MAX] = a[i:MAX] + b[i:MAX];
```

[Appendix] Preliminaries: Graphs



- Graph G = (V, E) consists of
 - Set of vertices V
 - Set of edges $E \subseteq V \times V$
- Edge $e = (u, v) \in E$: connection from source *u* to target *v*
- Number of vertices n = |V|
- Number of edges m = |E|
- Undirected graphs only: $(u, v) \in E$ iff $(v, u) \in E$
- Neighborhood $N(u) = \{v : (u, v) \in E\}$
- Loop $(u, u) \in E$
- Degree deg(v): number of incident edges, counting loops twice
- Distance: number of edges in shortest path connecting vertices
- Diameter d: greatest distance between any pair of vertices



[Appendix] Network Generation: Results



OpenMP Scheduling:



 \Rightarrow Scatter and balanced 1.4 times faster